

## ON A VARIANT OF PILLAI'S PROBLEM

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ABSTRACT. In this paper, we find all integers  $c$  having at least two representations as a difference between a Fibonacci number and a Tribonacci number.

## 1. INTRODUCTION

Pillai's famous conjecture first formulated in [11] states that the Diophantine equation

$$(1) \quad a^x - b^y = c$$

has for any fixed integer  $c > 0$  at most finitely many solutions  $a, b, x, y$  in positive integers. This conjecture is still open for all  $c \neq 1$ . Note that the case  $c = 1$  is Catalan's conjecture which has been solved by Mihăilescu [10]. If we leave  $a, b$  and  $c$  fixed, then much more is known about the solutions  $(x, y)$ . For instance Pillai [11] showed that if  $c$  is larger than some constant depending on  $a$  and  $b$ , then Diophantine equation (1) has at most one solution. In particular, he conjectured that in the case that  $a = 3$  and  $b = 2$  Diophantine equation (1) has at most one solution if  $c > 13$ . This conjecture was confirmed by Stroeker and Tijdeman [12] and their result was further improved by Bennett [5], who showed that for fixed  $a, b$  and  $c$  equation (1) has at most two solutions.

Recently Ddamulira, Luca and Rakotomalala [7] considered the Diophantine equation

$$(2) \quad F_n - 2^m = c,$$

where  $c$  is a fixed integer and  $\{F_n\}_{n \geq 0}$  is the sequence of Fibonacci numbers given by  $F_0 = 0$ ,  $F_1 = 1$  and  $F_{n+2} = F_{n+1} + F_n$  for all  $n \geq 0$ . This type of equation can be seen as a variation of Pillai's equation. However Ddamulira et.al. proved that the only integers  $c$  having at least two representations of the form  $F_n - 2^m$  are contained in the set  $\mathcal{C} = \{0, 1, -3, 5, -11, -30, 85\}$ . Moreover, they computed for all  $c \in \mathcal{C}$  all representations of the form (2).

The purpose of this paper is to consider a related problem. Denote by  $\{T_m\}_{m \geq 0}$  the sequence of Tribonacci numbers given by  $T_0 = 0$ ,  $T_1 = T_2 = 1$  and  $T_{m+3} = T_{m+2} + T_{m+1} + T_m$  for all  $m \geq 0$ . The main result of our paper is to find all nonzero integers  $c$  admitting at least two representations of the form  $F_n - T_m$  for some positive integers  $n$  and  $m$ . It is assumed that representations with  $n \in \{1, 2\}$  (for which  $F_1 = F_2 = 1$ ) as well as representations with  $m \in \{1, 2\}$  (for which

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$T_1 = T_2$ ) count as one representation to avoid trivial parametric families such as  $1 - 1 = F_1 - T_1 = F_2 - T_1 = F_1 - T_2 = F_2 - T_2$ . Therefore we assume that  $n \geq 2$  and  $m \geq 2$ . We prove the following theorem:

**Theorem 1.** *The only integers  $c$  having at least two representations of the form  $F_n - T_m$  come from the set*

$$\mathcal{C} = \{0, 1, -1, -2, -3, 4, -5, 6, 8, -10, 11, -11, -22, -23, -41, -60, -271\}.$$

*Furthermore, for each  $c \in \mathcal{C}$  all representations of the form  $c = F_n - T_m$  with integers  $n \geq 2$  and  $m \geq 2$  are:*

$$\begin{aligned} 0 &= 1 - 1 = 2 - 2 = 13 - 13 \quad (= F_2 - T_2 = F_3 - T_3 = F_7 - T_6), \\ 1 &= 2 - 1 = 3 - 2 = 5 - 4 = 8 - 7 \\ &\quad (= F_3 - T_2 = F_4 - T_3 = F_5 - T_4 = F_6 - T_5), \\ -1 &= 1 - 2 = 3 - 4 \quad (= F_2 - T_3 = F_4 - T_4), \\ -2 &= 2 - 4 = 5 - 7 \quad (= F_3 - T_4 = F_5 - T_5), \\ -3 &= 1 - 4 = 21 - 24 \quad (= F_2 - T_4 = F_8 - T_7), \\ 4 &= 5 - 1 = 8 - 4 \quad (= F_5 - T_2 = F_6 - T_4), \\ -5 &= 2 - 7 = 8 - 13 = 144 - 149 \quad (= F_3 - T_5 = F_6 - T_6 = F_{12} - T_{10}), \\ 6 &= 8 - 2 = 13 - 7 \quad (= F_6 - T_3 = F_7 - T_5), \\ 8 &= 21 - 13 = 89 - 81 \quad (= F_8 - T_6 = F_{11} - T_9), \\ -10 &= 3 - 13 = 34 - 44 \quad (= F_4 - T_6 = F_9 - T_8), \\ 11 &= 13 - 2 = 55 - 44 \quad (= F_7 - T_3 = F_{10} - T_8), \\ -11 &= 2 - 13 = 13 - 24 \quad (= F_3 - T_6 = F_7 - T_7), \\ -22 &= 2 - 24 = 121393 - 121415 \quad (= F_3 - T_7 = F_{26} - T_{21}), \\ -23 &= 1 - 24 = 21 - 44 \quad (= F_2 - T_7 = F_8 - T_8), \\ -41 &= 3 - 44 = 233 - 274 \quad (= F_4 - T_8 = F_{13} - T_{11}), \\ -60 &= 21 - 81 = 89 - 149 \quad (= F_8 - T_9 = F_{11} - T_{10}), \\ -271 &= 3 - 274 = 233 - 504 \quad (= F_4 - T_{11} = F_{13} - T_{12}). \end{aligned}$$

## 2. PRELIMINARIES

In this section, the result of linear forms in logarithms by Baker and Wüstholz [3] is stated. Besides, we state a lemma used by Ddamulira et.al. [7], which is a slight variation of a result due to Dujella and Pethő [8], of which is a generalization of a result due to Baker and Davenport [2]. Both results will be used in the proof of Theorem 1.

### 2.1. A lower bound for linear forms in logarithms of algebraic numbers.

In 1993, Baker and Wüstholz [3] obtained an explicit bound for linear forms in logarithms with a linear dependence on  $\log B$ . It is a vast improvement compared with lower bounds with a dependence on higher powers of  $\log B$  in preceding publications by other mathematicians in particular Baker's original results [1]. The final structure for the lower bound for linear forms in logarithms without an explicit determination of the constant  $C(k, d)$  involved has been established by Wüstholz [13] and the precise determination of that constant is the central aspect of [3] (see

also [4]). The improvement was mainly due to the use of the analytic subgroup theorem established by Wüstholz [14]. We shall now state the result of Baker and Wüstholz.

Denote by  $\alpha_1, \dots, \alpha_k$  algebraic numbers, not 0 or 1, and by  $\log \alpha_1, \dots, \log \alpha_k$  a fixed determination of their logarithms. Let  $K = \mathbb{Q}(\alpha_1, \dots, \alpha_k)$  and let  $d = [K : \mathbb{Q}]$  be the degree of  $K$  over  $\mathbb{Q}$ . For any  $\alpha \in K$ , suppose that its minimal polynomial over the integers is

$$g(x) = a_0 x^\delta + a_1 x^{\delta-1} + \dots + a_\delta = a_0 \prod_{j=1}^{\delta} (x - \alpha^{(j)}).$$

The absolute logarithmic Weil height of  $\alpha$  is defined as

$$h_0(\alpha) = \frac{1}{\delta} \left( \log |a_0| + \sum_{j=1}^{\delta} \log \left( \max\{|\alpha^{(j)}|, 1\} \right) \right).$$

Then the modified height  $h'(\alpha)$  is defined by

$$h'(\alpha) = \frac{1}{d} \max\{h(\alpha), |\log \alpha|, 1\},$$

where  $h(\alpha) = dh_0(\alpha)$  is the standard logarithmic Weil height of  $\alpha$ .

Let us consider the linear form

$$L(z_1, \dots, z_k) = b_1 z_1 + \dots + b_k z_k,$$

where  $b_1, \dots, b_k$  are rational integers, not all 0 and define

$$h'(L) = \frac{1}{d} \max\{h(L), 1\},$$

where  $h(L) = d \log \left( \max_{1 \leq j \leq k} \left\{ \frac{|b_j|}{b} \right\} \right)$  is the logarithmic Weil height of  $L$ , where  $b$  is the greatest common divisor of  $b_1, \dots, b_k$ . If we write  $B = \max\{|b_1|, \dots, |b_k|, e\}$ , then we get

$$h'(L) \leq \log B.$$

With these notations we are able to state the following result due to Baker and Wüstholz [3].

**Theorem 2.** *If  $\Lambda = L(\log \alpha_1, \dots, \log \alpha_k) \neq 0$ , then*

$$\log |\Lambda| \geq -C(k, d) h'(\alpha_1) \dots h'(\alpha_k) h'(L),$$

where

$$C(k, d) = 18(k+1)! k^{k+1} (32d)^{k+2} \log(2kd).$$

With  $|\Lambda| \leq \frac{1}{2}$ , we have  $\frac{1}{2} |\Lambda| \leq |\Phi| \leq 2|\Lambda|$ , where

$$\Phi = e^\Lambda - 1 = \alpha_1^{b_1} \dots \alpha_k^{b_k} - 1,$$

so that

$$\log |\alpha_1^{b_1} \dots \alpha_k^{b_k} - 1| \geq \log |\Lambda| - \log 2.$$

We apply Theorem 2 mainly in a situation where  $k = 3$  and  $d = 6$ . In this case we obtain

$$C(3, 6) = 18(4!) 3^4 (32 \times 6)^5 (\log 36) \approx 3.2718 \dots \times 10^{16}.$$

We will use this value throughout the paper without any further reference.

**2.2. A generalized result by Dujella and Pethő.** The following result will be used to reduce the huge upper bounds for  $n$  and  $m$  which appear during the course of the proof of Theorem 1 (cf. Proposition 1). The following Lemma is stated in [7], which is regarded as a slight variation of a result due to Dujella and Pethő [8], of which is a generalization of a result due to Baker and Davenport [2]. For a real number  $x$ , let  $\|x\| = \min\{|x - n| : n \in \mathbb{Z}\}$  be the distance from  $x$  to the nearest integer.

**Lemma 1.** *Let  $M$  be a positive integer, let  $p/q$  be a convergent of the continued fraction of the irrational  $\tau$  such that  $q > 6M$ , and let  $A, B, \mu$  be some real numbers with  $A > 0$  and  $B > 1$ . Let  $\varepsilon := \|\mu q\| - M\|\tau q\|$ . If  $\varepsilon > 0$ , then there is no solution to the inequality*

$$0 < m\tau - n + \mu < AB^{-k},$$

*in positive integers  $m, n$  and  $k$  with*

$$m \leq M \quad \text{and} \quad k \geq \frac{\log(Aq/\varepsilon)}{\log B}.$$

### 3. PROOF OF THEOREM 1

**3.1. Set up.** Assume that  $(n, m) \neq (n_1, m_1)$  are pairs of indices such that

$$(3) \quad F_n - F_{n_1} = T_m - T_{m_1}.$$

We may assume that  $m \neq m_1$ , since otherwise  $(n, m) = (n_1, m_1)$ . Furthermore we assume that  $m > m_1$ . Due to equation (3) and since the right hand side of equation (3) is positive, we get that the left hand side of equation (3) is also positive and thus  $n > n_1$ . Therefore, we have  $n \geq 3$ ,  $n_1 \geq 2$  and  $m \geq 3$ ,  $m_1 \geq 2$ .

During the proof of Theorem 1 we use the Binet formulae for the Fibonacci sequence and Tribonacci sequence in the following form:

**Fibonacci sequence:**

$$F_k = \frac{\alpha^k - \beta^k}{\alpha - \beta} \quad \text{for all } k \geq 0,$$

where  $\alpha = \frac{1+\sqrt{5}}{2}$  and  $\beta = \frac{1-\sqrt{5}}{2}$  are the roots of the characteristic equation  $x^2 - x - 1 = 0$ . Besides, the inequality

$$\alpha^{k-2} \leq F_k \leq \alpha^{k-1}$$

holds for all  $k \geq 1$ .

**Tribonacci sequence:**

$$T_k = c_\alpha \alpha_T^k + c_\beta \beta_T^k + c_\gamma \gamma_T^k \quad \text{for all } k \geq 0,$$

where  $\alpha_T$ ,  $\beta_T$  and  $\gamma_T$  are the roots of the characteristic equation  $x^3 - x^2 - x - 1 = 0$ , with

$$\begin{aligned} \alpha_T &= \frac{1}{3} \left( 1 + \sqrt[3]{19 + 3\sqrt{33}} + \sqrt[3]{19 - 3\sqrt{33}} \right), \\ \beta_T &= \frac{1}{6} \left( 2 - \sqrt[3]{19 + 3\sqrt{33}} - \sqrt[3]{19 - 3\sqrt{33}} \right) + \frac{\sqrt{3}}{6} i \left( \sqrt[3]{19 + 3\sqrt{33}} - \sqrt[3]{19 - 3\sqrt{33}} \right), \\ \gamma_T &= \frac{1}{6} \left( 2 - \sqrt[3]{19 + 3\sqrt{33}} - \sqrt[3]{19 - 3\sqrt{33}} \right) - \frac{\sqrt{3}}{6} i \left( \sqrt[3]{19 + 3\sqrt{33}} - \sqrt[3]{19 - 3\sqrt{33}} \right), \end{aligned}$$

and the coefficients

$$\begin{aligned} c_\alpha &= \frac{\alpha_T}{(\alpha_T - \beta_T)(\alpha_T - \gamma_T)} = \frac{1}{-\alpha_T^2 + 4\alpha_T - 1}, \\ c_\beta &= \frac{\beta_T}{(\beta_T - \alpha_T)(\beta_T - \gamma_T)} = \frac{1}{-\beta_T^2 + 4\beta_T - 1}, \\ c_\gamma &= \frac{\gamma_T}{(\gamma_T - \alpha_T)(\gamma_T - \beta_T)} = \frac{1}{-\gamma_T^2 + 4\gamma_T - 1} \end{aligned}$$

are the roots of the polynomial  $44x^3 - 2x - 1$ . Note that

$$\begin{aligned} 1.839 < \alpha_T < 1.840 & \quad 0.336 < c_\alpha < 0.337 \\ \beta_T = \overline{\gamma_T} & \quad 0.737 < |\beta_T| = |\gamma_T| < 0.738 \\ c_\beta = \overline{c_\gamma} & \quad 0.259 < |c_\beta| = |c_\gamma| < 0.260. \end{aligned}$$

Finally let us state several useful inequalities. For instance

$$\alpha_T^{k-2} \leq T_k \leq \alpha_T^{k-1} \quad \text{for all } k \geq 1.$$

which was already shown in [6]. Using equation (3) we get that

$$(4) \quad \alpha^{n-4} \leq F_{n-2} \leq F_n - F_{n_1} = T_m - T_{m_1} < T_m \leq \alpha_T^{m-1},$$

and similarly we get

$$(5) \quad \begin{aligned} \alpha^{n-1} &\geq F_n > F_n - F_{n_1} = T_m - T_{m_1} \geq T_m - T_{m-1} = \\ &T_{m-2} + T_{m-3} \geq \alpha_T^{m-4} + \alpha_T^{m-5} > 2.83\alpha_T^{m-5}. \end{aligned}$$

Thus

$$(6) \quad n - 4 < \frac{\log \alpha_T}{\log \alpha}(m - 1) \quad \text{and} \quad n - 3 > \frac{\log \alpha_T}{\log \alpha}(m - 5),$$

where  $\frac{\log \alpha_T}{\log \alpha} \approx 1.2663 \dots$

Inequality (6) implies that if  $n < 300$ , then  $m < 240$ . By a brute force computer enumeration for  $2 \leq n_1 < n < 300$  and  $2 \leq m_1 < m < 240$  we found all solutions listed in Theorem 1. Thus we may assume from now on that  $n \geq 300$ .

**3.2. Linear forms in logarithms.** Since  $n \geq 300$ , by the first inequality of (6) we obtain that  $m \geq 235$  which combined with the second inequality of (6) implies that  $n > m$ . Moreover, we have

$$\frac{\alpha^n - \beta^n}{\sqrt{5}} - \frac{\alpha^{n_1} - \beta^{n_1}}{\sqrt{5}} = (c_\alpha \alpha_T^m + c_\beta \beta_T^m + c_\gamma \gamma_T^m) - (c_\alpha \alpha_T^{m_1} + c_\beta \beta_T^{m_1} + c_\gamma \gamma_T^{m_1}).$$

Collecting the “large” terms on the left hand side of the equation we obtain

$$\begin{aligned} \left| \frac{\alpha^n}{\sqrt{5}} - c_\alpha \alpha_T^m \right| &= \left| \frac{\beta^n}{\sqrt{5}} + \frac{\alpha^{n_1} - \beta^{n_1}}{\sqrt{5}} + (c_\beta \beta_T^m + c_\gamma \gamma_T^m) - (c_\alpha \alpha_T^{m_1} + c_\beta \beta_T^{m_1} + c_\gamma \gamma_T^{m_1}) \right| \\ &\leq \frac{\alpha^{n_1}}{\sqrt{5}} + c_\alpha \alpha_T^{m_1} + \frac{|\beta|^n}{\sqrt{5}} + \frac{|\beta|^{n_1}}{\sqrt{5}} + |c_\beta| |\beta_T|^m + |c_\gamma| |\gamma_T|^m \\ &\quad + |c_\beta| |\beta_T|^{m_1} + |c_\gamma| |\gamma_T|^{m_1} \\ &< \frac{\alpha^{n_1}}{\sqrt{5}} + c_\alpha \alpha_T^{m_1} + 0.46 \\ &< 0.92 \max\{\alpha^{n_1}, \alpha_T^{m_1}\}. \end{aligned}$$

Dividing by  $c_\alpha \alpha_T^m$  we get

$$\begin{aligned} \left| (\sqrt{5}c_\alpha)^{-1} \alpha^n \alpha_T^{-m} - 1 \right| &< \max \left\{ \frac{0.92}{c_\alpha \alpha_T^m} \alpha^{n_1}, \frac{0.92}{c_\alpha} \alpha_T^{m_1-m} \right\} \\ &< \max \left\{ 2.74 \frac{\alpha^{n_1}}{\alpha_T} \frac{1}{\alpha^{n-4}}, 2.74 \alpha_T^{m_1-m} \right\}. \end{aligned}$$

Hence we obtain the inequality

$$(7) \quad \left| (\sqrt{5}c_\alpha)^{-1} \alpha^n \alpha_T^{-m} - 1 \right| < \max \{ \alpha^{n_1-n+5}, \alpha_T^{m_1-m+2} \}.$$

Let us introduce

$$\Lambda = n \log \alpha - m \log \alpha_T - \log(\sqrt{5}c_\alpha)$$

and assume that  $|\Lambda| \leq 0.5$ . Further, we put

$$\Phi = e^\Lambda - 1 = (\sqrt{5}c_\alpha)^{-1} \alpha^n \alpha_T^{-m} - 1$$

and use the theorem of Baker and Wüstholz (Theorem 2) with the data

$$k = 3, \quad \alpha_1 = \sqrt{5}c_\alpha, \quad b_1 = -1, \quad \alpha_2 = \alpha, \quad b_2 = n, \quad \alpha_3 = \alpha_T, \quad b_3 = -m.$$

With this data we have  $K = \mathbb{Q}(\sqrt{5}, \alpha_T)$ , i.e.  $d = 6$ , and  $B = n$ . Notice that the minimal polynomial of  $\alpha_1$  is  $1936x^6 - 880x^4 + 100x^2 - 125$ , and we conclude that  $h'(\alpha_1) = \frac{1}{6} \log 1936$ . Further we obtain by a simple computation that  $h'(\alpha_2) = \frac{1}{2} \log \alpha$  and  $h'(\alpha_3) = \frac{1}{3} \log \alpha_T$ .

Before we can apply Theorem 2 we have to show that  $\Phi \neq 0$ . Assume to the contrary that  $\Phi = 0$ , then  $\alpha^{2n} = 5c_\alpha^2 \alpha_T^{2m}$ . But  $\alpha^{2n} \in \mathbb{Q}(\sqrt{5}) \setminus \mathbb{Q}$  whereas  $5c_\alpha^2 \alpha_T^{2m} \in \mathbb{Q}(\alpha_T)$ . Thus  $\Phi = 0$  is impossible due to the fact that  $\mathbb{Q}(\sqrt{5}) \cap \mathbb{Q}(\alpha_T) = \mathbb{Q}$ .

Apply Theorem 2 yields

$$\log |\Phi| \geq -C(3, 6) \left( \frac{1}{6} \log 1936 \right) \left( \frac{1}{2} \log \alpha \right) \left( \frac{1}{3} \log \alpha_T \right) \log n - \log 2$$

and together with inequality (7) we have

$$\min \{ (n - n_1 - 5) \log \alpha, (m - m_1 - 2) \log \alpha_T \} \leq 2.02 \times 10^{15} \log n.$$

Thus we have proved so far:

**Lemma 2.** *Assume that  $(n, m, n_1, m_1)$  is a solution to equation (3) with  $m > m_1$ . Then we have*

$$\min \{ (n - n_1) \log \alpha, (m - m_1) \log \alpha_T \} < 2.03 \times 10^{15} \log n.$$

Note that in the case that  $|\Lambda| > 0.5$  inequality (7) is possible only if either  $n - n_1 \leq 5$  or  $m - m_1 \leq 2$ , which is covered by the bound provided by Lemma 2.

Now we have to distinguish between the following two cases:

**Case 1.** Let us assume that

$$\min \{ (n - n_1) \log \alpha, (m - m_1) \log \alpha_T \} = (n - n_1) \log \alpha.$$

We rewrite equation (3) as

$$\begin{aligned} \left| \frac{\alpha^n - \alpha^{n_1}}{\sqrt{5}} - c_\alpha \alpha_T^m \right| &= \left| -c_\alpha \alpha_T^{m_1} + \frac{\beta^n}{\sqrt{5}} - \frac{\beta^{n_1}}{\sqrt{5}} + \right. \\ &\quad \left. (c_\beta \beta_T^m + c_\gamma \gamma_T^m) - (c_\beta \beta_T^{m_1} + c_\gamma \gamma_T^{m_1}) \right| \\ &\leq c_\alpha \alpha_T^{m_1} + \frac{|\beta^n|}{\sqrt{5}} + \frac{|\beta^{n_1}|}{\sqrt{5}} + |c_\beta| |\beta_T^m| + |c_\gamma| |\gamma_T^m| + \\ &\quad |c_\beta| |\beta_T^{m_1}| + |c_\gamma| |\gamma_T^{m_1}| \end{aligned}$$

and obtain that

$$\left| \frac{\alpha^{n-n_1} - 1}{\sqrt{5}} \alpha^{n_1} - c_\alpha \alpha_T^m \right| < (c_\alpha + 0.14) \alpha_T^{m_1}.$$

Dividing by  $c_\alpha \alpha_T^m$  we get the inequality

$$(8) \quad \left| \frac{\alpha^{n-n_1} - 1}{\sqrt{5} c_\alpha} \alpha^{n_1} \alpha_T^{-m} - 1 \right| < 1.42 \alpha_T^{m_1-m}.$$

**Case 2.** Let us assume that

$$\min\{(n - n_1) \log \alpha, (m - m_1) \log \alpha_T\} = (m - m_1) \log \alpha_T.$$

We rewrite equation (3) as

$$\begin{aligned} \left| \frac{\alpha^n}{\sqrt{5}} - c_\alpha \alpha_T^m + c_\alpha \alpha_T^{m_1} \right| &= \left| \frac{\beta^n}{\sqrt{5}} + \frac{\alpha^{n_1} - \beta^{n_1}}{\sqrt{5}} + c_\beta \beta_T^m + c_\gamma \gamma_T^m - c_\beta \beta_T^{m_1} - c_\gamma \gamma_T^{m_1} \right| \\ &\leq \frac{|\beta^n|}{\sqrt{5}} + \frac{\alpha^{n_1}}{\sqrt{5}} + \frac{|\beta^{n_1}|}{\sqrt{5}} + |c_\beta| |\beta_T^m| + |c_\gamma| |\gamma_T^m| \\ &\quad + |c_\beta| |\beta_T^{m_1}| + |c_\gamma| |\gamma_T^{m_1}|. \end{aligned}$$

Thus we get

$$\left| \alpha^n - \sqrt{5} c_\alpha (\alpha_T^{m-m_1} - 1) \alpha_T^{m_1} \right| < 1.4 \alpha^{n_1}.$$

Dividing both sides by  $\sqrt{5} c_\alpha (\alpha_T^{m-m_1} - 1) \alpha_T^{m_1}$  we get by using inequality (4) the following inequality:

$$(9) \quad \left| \frac{\alpha^n \alpha_T^{-m_1}}{\sqrt{5} c_\alpha (\alpha_T^{m-m_1} - 1)} - 1 \right| < \frac{1.4}{\sqrt{5} c_\alpha (1 - \alpha_T^{m_1-m}) \alpha_T} \frac{\alpha^{n_1}}{\alpha_T^{m-1}} < 2.22 \alpha^{n_1-n+4}.$$

We want to apply Theorem 2 to both inequalities (8) and (9) respectively. Let us consider the first case more closely. We write

$$\Lambda_1 = n_1 \log \alpha - m \log \alpha_T + \log \left( \frac{\alpha^{n-n_1} - 1}{\sqrt{5} c_\alpha} \right)$$

and assume that  $|\Lambda_1| \leq 0.5$ . Further, we put

$$\Phi_1 = e^{\Lambda_1} - 1 = \frac{\alpha^{n-n_1} - 1}{\sqrt{5} c_\alpha} \alpha^{n_1} \alpha_T^{-m} - 1$$

and aim to apply Theorem 2 by taking  $K = \mathbb{Q}(\sqrt{5}, \alpha_T)$ , i.e.  $d = 6$ ,  $k = 3$  and  $B = n$ . Further, we have

$$\alpha_1 = \frac{\alpha^{n-n_1} - 1}{\sqrt{5} c_\alpha}, \quad b_1 = 1, \quad \alpha_2 = \alpha, \quad b_2 = n_1, \quad \alpha_3 = \alpha_T, \quad b_3 = -m.$$

Let us estimate the height of  $\alpha_1$ . Notice that  $h(\alpha_1) \leq h(\eta_1) + h(\eta_2)$ , where  $\eta_1 = \frac{\alpha^{n-n_1}-1}{\sqrt{5}}$  and  $\eta_2 = \frac{1}{c_\alpha}$ . The minimal polynomial of  $\eta_1$  divides (e.g. see [7])

$$5X^2 - 5F_{n-n_1}X - ((-1)^{n-n_1} + 1 - L_{n-n_1}),$$

where  $\{L_k\}_{k \geq 0}$  is the Lucas companion sequence of the Fibonacci sequence given by  $L_0 = 2, L_1 = 1, L_{k+2} = L_{k+1} + L_k$  for  $k \geq 0$ . Its Binet formula for the general term is  $L_k = \alpha^k + \beta^k$  for all  $k \geq 0$ . Thus (cf. [7]),

$$h_0(\eta_1) \leq \frac{1}{2} \left( \log 5 + \log \left( \frac{\alpha^{n-n_1} + 1}{\sqrt{5}} \right) \right).$$

Thus Lemma 2 yields an upper bound

$$h_0(\eta_1) < \frac{1}{2} \log (2\sqrt{5}\alpha^{n-n_1}) < \frac{1}{2}(n - n_1 + 4) \log \alpha < 1.02 \times 10^{15} \log n,$$

i.e.  $h(\eta_1) < 6 \times 1.02 \times 10^{15} \log n$ . Since  $h_0(\eta_2) = h_0(c_\alpha) = \frac{1}{3} \log 44$ , i.e.  $h(\eta_2) = 2 \log 44$ , we have  $h(\alpha_1) \leq 6 \times 1.02 \times 10^{15} \log n + 2 \log 44$  and finally we obtain that

$$h'(\alpha_1) < 1.03 \times 10^{15} \log n.$$

Moreover, we have that  $h'(\alpha_2) = \frac{1}{2} \log \alpha$  and  $h'(\alpha_3) = \frac{1}{3} \log \alpha_T$  as before.

Now let us turn to the second case. We write

$$A_2 = n \log \alpha - m_1 \log \alpha_T - \log \left( \sqrt{5} c_\alpha (\alpha_T^{m-m_1} - 1) \right)$$

and assume that  $|A_2| \leq 0.5$ . Further, we put

$$\Phi_2 = e^{A_2} - 1 = (\sqrt{5} c_\alpha (\alpha_T^{m-m_1} - 1))^{-1} \alpha^n \alpha_T^{-m_1} - 1$$

and aim to apply Theorem 2. As in the previous case we take  $K = \mathbb{Q}(\sqrt{5}, \alpha_T)$ , i.e.  $d = 6$ ,  $k = 3$  and  $B = n$ . Further, we have

$$\alpha_1 = \sqrt{5} c_\alpha (\alpha_T^{m-m_1} - 1), \quad b_1 = -1, \quad \alpha_2 = \alpha, \quad b_2 = n, \quad \alpha_3 = \alpha_T, \quad b_3 = -m_1.$$

Again, we have to estimate  $h(\alpha_1)$  and therefore note that  $h(\alpha_1) \leq h(\eta_1) + h(\eta_2) + h(\eta_3)$ , where  $\eta_1 = \alpha_T^{m-m_1} - 1$ ,  $\eta_2 = c_\alpha$  and  $\eta_3 = \sqrt{5}$ . By applying Lemma 2 we get

$$\begin{aligned} h_0(\eta_1) &\leq h_0(\alpha_T^{m-m_1}) + h_0(-1) + \log 2 \\ &= (m - m_1) h_0(\alpha_T) + \log 2 = \frac{m - m_1}{3} \log \alpha_T + \log 2 \\ &< \frac{1}{3} \times 2.03 \times 10^{15} \log n + \log 2. \end{aligned}$$

Thus

$$h(\alpha_1) < 6 \left( \frac{1}{3} \times 2.03 \times 10^{15} \log n + \log 2 + \frac{1}{3} \log 44 + \log \sqrt{5} \right)$$

and therefore

$$h'(\alpha_1) < 6.77 \times 10^{14} \log n < 1.03 \times 10^{15} \log n.$$

Once again, we have that  $h'(\alpha_2) = \frac{1}{2} \log \alpha$  and  $h'(\alpha_3) = \frac{1}{3} \log \alpha_T$ .

In particular, we have shown in both cases that

$$h'(\alpha_1) < 1.03 \times 10^{15} \log n, \quad h'(\alpha_2) = \frac{1}{2} \log \alpha, \quad h'(\alpha_3) = \frac{1}{3} \log \alpha_T, \quad B = n.$$

But, before we can apply Theorem 2 we have to ensure that  $\Phi_i \neq 0$  for  $i = 1, 2$ . Firstly we deal with the assumption that  $\Phi_1 = 0$ , i.e.  $\alpha^n - \alpha^{n_1} = \sqrt{5} c_\alpha \alpha_T^m$ . This is impossible if  $\sqrt{5} c_\alpha \alpha_T^m \in \mathbb{Q}(\sqrt{5}, \alpha_T)$  but  $\notin \mathbb{Q}(\sqrt{5})$ . Therefore let us assume that



$\sqrt{5}c_\alpha\alpha_T^m \in \mathbb{Q}(\sqrt{5})$ , hence  $\sqrt{5}c_\alpha\alpha_T^m = y\sqrt{5}$  for some  $y \in \mathbb{Q}$ . Let  $\sigma \neq \text{id}$  be the unique non-trivial  $\mathbb{Q}$ -automorphism over  $\mathbb{Q}(\sqrt{5})$ . Then we get

$$\alpha^n - \alpha^{n_1} = \sqrt{5}c_\alpha\alpha_T^m = y\sqrt{5} = -\sigma(\sqrt{5}c_\alpha\alpha_T^m) = -\sigma(\alpha^n - \alpha^{n_1}) = \beta^{n_1} - \beta^n.$$

However, the absolute value of  $\alpha^n - \alpha^{n_1}$  is at least  $\alpha^n - \alpha^{n_1} \geq \alpha^{n-2} \geq \alpha^{298} > 2$  whereas the absolute value of  $\beta^{n_1} - \beta^n$  is at most  $|\beta^{n_1} - \beta^n| \leq |\beta|^{n_1} + |\beta|^n < 2$ . By this obvious contradiction we conclude that  $\Phi_1 \neq 0$ .

Now let us consider the second case and assume for the moment that  $\Phi_2 = 0$ , i.e.  $\alpha^{2n} = 5\alpha_T^{2m_1}c_\alpha^2(\alpha_T^{m-m_1} - 1)^2$ . However,  $\alpha^{2n} \in \mathbb{Q}(\sqrt{5}) \setminus \mathbb{Q}$ , whereas  $5\alpha_T^{2m_1}c_\alpha^2(\alpha_T^{m-m_1} - 1)^2 \in \mathbb{Q}(\alpha_T)$ . Thus we obtain also in this case a contradiction and we also conclude in this case that  $\Phi_2 \neq 0$ .

Now, we are ready to apply Theorem 2 and get

$$\begin{aligned} \log |\Phi_i| &> -C(3, 6) (1.03 \times 10^{15} \log n) \left( \frac{1}{2} \log \alpha \right) \left( \frac{1}{3} \log \alpha_T \right) \log n - \log 2 \\ &> -1.65 \times 10^{30} (\log n)^2 \end{aligned}$$

for  $i = 1, 2$ . Combining this inequality with the inequalities (8) and (9), we obtain

$$(m_1 - m) \log \alpha_T + \log 1.42 > -1.65 \times 10^{30} (\log n)^2$$

and

$$(n_1 - n + 4) \log \alpha + \log 2.22 > -1.65 \times 10^{30} (\log n)^2$$

respectively. These two inequalities yield together with Lemma 2 the following lemma:

**Lemma 3.** *Assume that  $(n, m, n_1, m_1)$  is a solution to equation (3) with  $m > m_1$ . Then we have*

$$\max\{(n - n_1) \log \alpha, (m - m_1) \log \alpha_T\} < 1.66 \times 10^{30} (\log n)^2.$$

Note that in view of inequality (8)  $|A_1| > 0.5$  is possible only if  $m - m_1 = 1$  and in view of inequality (9)  $|A_2| > 0.5$  is possible only if  $n - n_1 \leq 6$  respectively. Both cases are covered by the bound provided by Lemma 3.

One more time we have to apply Theorem 2. This time we rewrite equation (3) as

$$\begin{aligned} \left| \frac{\alpha^n}{\sqrt{5}} - \frac{\alpha^{n_1}}{\sqrt{5}} - c_\alpha\alpha_T^m + c_\alpha\alpha_T^{m_1} \right| \\ = \left| \frac{\beta^n}{\sqrt{5}} - \frac{\beta^{n_1}}{\sqrt{5}} + c_\beta\beta_T^m + c_\gamma\gamma_T^m - c_\beta\beta_T^{m_1} - c_\gamma\gamma_T^{m_1} \right| < 0.46 \end{aligned}$$

Dividing both sides by  $c_\alpha\alpha_T^{m_1}(\alpha_T^{m-m_1} - 1)$  we get by applying inequality (4)

$$(10) \quad \left| \frac{\alpha^{n-n_1} - 1}{\sqrt{5}c_\alpha(\alpha_T^{m-m_1} - 1)} \alpha^{n_1} \alpha_T^{-m_1} - 1 \right| < \frac{0.46}{c_\alpha(1 - \alpha_T^{m_1-m})\alpha_T} \frac{1}{\alpha_T^{m-1}} < 1.64\alpha^{4-n}.$$

In this final step we consider the linear form

$$A_3 = n_1 \log \alpha - m_1 \log \alpha_T + \log \left( \frac{\alpha^{n-n_1} - 1}{\sqrt{5}c_\alpha(\alpha_T^{m-m_1} - 1)} \right)$$

and assume that  $|\Lambda_3| \leq 0.5$ . Further, we put

$$\Phi_3 = e^{\Lambda_3} - 1 = \frac{\alpha^{n-n_1} - 1}{\sqrt{5} c_\alpha (\alpha_T^{m-m_1} - 1)} \alpha^{n_1} \alpha_T^{-m_1} - 1.$$

As before we take  $K = \mathbb{Q}(\sqrt{5}, \alpha_T)$ , i.e.  $d = 6$ ,  $k = 3$ ,  $B = n$  and we have

$$\alpha_1 = \frac{\alpha^{n-n_1} - 1}{\sqrt{5} c_\alpha (\alpha_T^{m-m_1} - 1)}, \quad b_1 = 1, \quad \alpha_2 = \alpha, \quad b_2 = n_1, \quad \alpha_3 = \alpha_T, \quad b_3 = -m_1.$$

By Lemma 3 and similar computations as done before we obtain that

$$h\left(\frac{\alpha^{n-n_1} - 1}{\sqrt{5} c_\alpha}\right) \leq 6 \left(\frac{1}{2}(n - n_1 + 4) \log \alpha\right) + 2 \log 44 < 3 \times (1.67 \times 10^{30} (\log n)^2)$$

and

$$h(\alpha_T^{m-m_1} - 1) \leq 6 \left(\frac{m - m_1}{3} \log \alpha_T + \log 2\right) < 2 \times (1.67 \times 10^{30} (\log n)^2).$$

Therefore we find the upper bound

$$h(\alpha_1) \leq h\left(\frac{\alpha^{n-n_1} - 1}{\sqrt{5} c_\alpha}\right) + h(\alpha_T^{m-m_1} - 1) < 5 \times (1.67 \times 10^{30} (\log n)^2)$$

and thus

$$h'(\alpha_1) < \frac{5}{6} \times (1.67 \times 10^{30} (\log n)^2).$$

As before, we have  $h'(\alpha_2) = \frac{1}{2} \log \alpha$  and  $h'(\alpha_3) = \frac{1}{3} \log \alpha_T$ .

Using similar arguments as in the proof that  $\Phi_1 \neq 0$  we can show that  $\Phi_3 \neq 0$ . Now an application of Theorem 2 yields

$$\log |\Phi_3| > -C(3, 6) \left(\frac{5}{6} \times 1.67 \times 10^{30} (\log n)^2\right) \left(\frac{1}{2} \log \alpha\right) \left(\frac{1}{3} \log \alpha_T\right) \log n - \log 2.$$

Combining this inequality with inequality (10) we get

$$(n - 4) \log \alpha < 2.23 \times 10^{45} (\log n)^3$$

which yields  $n < 8 \times 10^{51}$ .

Similarly as in the cases above the assumption that  $|\Lambda_3| > 0.5$  leads in view of inequality (10) to  $n \leq 5$ . Let us summarize the results of this subsection:

**Proposition 1.** *Assume that  $(n, m, n_1, m_1)$  is a solution to equation (3) with  $m > m_1$ . Then we have that  $n < 8 \times 10^{51}$ .*

*Remark 1.* The theorem of Baker and Wüstholz (Theorem 2) can be easily applied. However, a slightly sharper bound for  $n$ , namely  $n < 6 \times 10^{48}$ , may be obtained if one uses Matveev's result [9] instead. However, the improvement is not crucial in view of our next step, the reduction of our upper bound for  $n$  using the method of Baker and Davenport (Lemma 1).

**3.3. Generalized method of Baker and Davenport.** In our final step we reduce the huge upper bound for  $n$  from Proposition 1 by applying several times Lemma 1. In this subsection we follow the ideas from [7]. First, we consider inequality (7) and recall that

$$\Lambda = n \log \alpha - m \log \alpha_T - \log(\sqrt{5}c_\alpha).$$

For technical reasons we assume that  $\min\{n - n_1, m - m_1\} \geq 20$ . In the case that this condition fails we consider one of the following inequalities instead:

- if  $n - n_1 < 20$  but  $m - m_1 \geq 20$ , we consider inequality (8);
- if  $n - n_1 \geq 20$  but  $m - m_1 < 20$ , we consider inequality (9);
- if both  $n - n_1 < 20$  and  $m - m_1 < 20$ , we consider inequality (10).

Let us start by considering inequality (7). Since we assume that  $\min\{n - n_1, m - m_1\} \geq 20$  we get  $|\Phi| = |e^\Lambda - 1| < \frac{1}{4}$ , hence  $|\Lambda| < \frac{1}{2}$ . And, since  $|x| < 2|e^x - 1|$  holds for all  $x \in (-\frac{1}{2}, \frac{1}{2})$  we get

$$|\Lambda| < 2 \max\{\alpha^{n_1-n+5}, \alpha_T^{m_1-m+2}\} \leq \max\{\alpha^{n_1-n+7}, \alpha_T^{m_1-m+4}\}.$$

Assume that  $\Lambda > 0$ . Then we have the inequality

$$\begin{aligned} 0 < n \left( \frac{\log \alpha}{\log \alpha_T} \right) - m + \frac{\log(1/(\sqrt{5}c_\alpha))}{\log \alpha_T} &< \max \left\{ \frac{\alpha^7}{\log \alpha_T} \alpha^{-(n-n_1)}, \frac{\alpha_T^4}{\log \alpha_T} \alpha_T^{-(m-m_1)} \right\} \\ &< \max \left\{ 48\alpha^{-(n-n_1)}, 19\alpha_T^{-(m-m_1)} \right\} \end{aligned}$$

and we apply Lemma 1 with

$$\tau = \frac{\log \alpha}{\log \alpha_T}, \quad \mu = \frac{\log(1/(\sqrt{5}c_\alpha))}{\log \alpha_T}, \quad (A, B) = (48, \alpha) \text{ or } (19, \alpha_T).$$

Let  $\tau = [a_0, a_1, a_2, \dots] = [0, 1, 3, 1, 3, 13, 2, 1, 8, 3, 1, 5, \dots]$  be the continued fraction of  $\tau$ . Moreover, we choose  $M = 8 \times 10^{51}$  and consider the 104-th convergent

$$\frac{p}{q} = \frac{p_{104}}{q_{104}} = \frac{528419636478855291192208008138409657842309076397924033}{669159011284129920139468279297504453112608160771671810},$$

with  $q = q_{104} > 6M$ . This yields  $\varepsilon > 0.068$  and therefore either

$$n - n_1 \leq \frac{\log(48q/0.068)}{\log \alpha} < 272, \quad \text{or} \quad m - m_1 \leq \frac{\log(19q/0.068)}{\log \alpha_T} < 213.$$

Thus, we have either  $n - n_1 \leq 271$ , or  $m - m_1 \leq 212$ .

In the case of  $\Lambda < 0$  we consider the following inequality:

$$\begin{aligned} 0 < m \left( \frac{\log \alpha_T}{\log \alpha} \right) - n + \frac{\log(\sqrt{5}c_\alpha)}{\log \alpha} &< \max \left\{ \frac{\alpha^7}{\log \alpha_T} \alpha^{-(n-n_1)}, \frac{\alpha_T^4}{\log \alpha_T} \alpha_T^{-(m-m_1)} \right\} \\ &< \max \left\{ 61\alpha^{-(n-n_1)}, 24\alpha_T^{-(m-m_1)} \right\} \end{aligned}$$

instead and apply Lemma 1 with

$$\tau = \frac{\log \alpha_T}{\log \alpha}, \quad \mu = \frac{\log(\sqrt{5}c_\alpha)}{\log \alpha}, \quad (A, B) = (61, \alpha) \text{ or } (24, \alpha_T).$$

Let  $\tau = [a_0, a_1, a_2, \dots] = [1, 3, 1, 3, 13, 2, 1, 8, 3, 1, 5, 2, \dots]$  be the continued fraction of  $\tau$ . Again, we choose  $M = 8 \times 10^{51}$  but in this case we consider instead of the 104-th convergent the 103-rd convergent

$$\frac{p}{q} = \frac{p_{103}}{q_{103}} = \frac{669159011284129920139468279297504453112608160771671810}{528419636478855291192208008138409657842309076397924033},$$

with  $q > 6M$ . This yields  $\varepsilon > 0.067$  and again we obtain either

$$n - n_1 \leq \frac{\log(61q/0.067)}{\log \alpha} < 272, \quad \text{or} \quad m - m_1 \leq \frac{\log(24q/0.067)}{\log \alpha_T} < 213.$$

These bounds agree with the bounds obtained in the case that  $\Lambda > 0$ . As a conclusion, we have either  $n - n_1 \leq 271$  or  $m - m_1 \leq 212$  whenever  $\Lambda \neq 0$ .

Now, we have to distinguish between the two cases  $n - n_1 \leq 271$  and  $m - m_1 \leq 212$ . First, let us assume that  $n - n_1 \leq 271$ . In this case we consider inequality (8) and assume that  $m - m_1 \geq 20$ . Recall that

$$\Lambda_1 = n_1 \log \alpha - m \log \alpha_T + \log \left( \frac{\alpha^{n-n_1} - 1}{\sqrt{5}c_\alpha} \right)$$

and inequality (8) yields that

$$|\Lambda_1| < \alpha_T^{m_1-m+2}.$$

If we further assume that  $\Lambda_1 > 0$ , then we get

$$0 < n_1 \left( \frac{\log \alpha}{\log \alpha_T} \right) - m + \frac{\log((\alpha^{n-n_1} - 1)/(\sqrt{5}c_\alpha))}{\log \alpha_T} < \frac{\alpha_T^2}{\log \alpha_T} \alpha_T^{-(m-m_1)} < 6\alpha_T^{-(m-m_1)}.$$

Again we apply Lemma 1 with the same  $\tau$  and  $M$  as in the case that  $\Lambda > 0$ . We use the 104-th convergent  $\frac{p}{q} = \frac{p_{104}}{q_{104}}$  of  $\tau$  as before. But, in this case we choose  $(A, B) = (6, \alpha_T)$  and use

$$\mu_k = \frac{\log((\alpha^k - 1)/(\sqrt{5}c_\alpha))}{\log \alpha_T},$$

instead of  $\mu$  for each possible value of  $n - n_1 = k = 1, 2, \dots, 271$ . A quick computer aid computation yields that  $\varepsilon > 0.00038$  for all  $1 \leq k \leq 271$ . Hence, by Lemma 1, we get

$$m - m_1 < \frac{\log(6q/0.00038)}{\log \alpha_T} < 220.$$

Thus,  $n - n_1 \leq 271$  implies  $m - m_1 \leq 219$ .

In the case that  $\Lambda_1 < 0$  we follow the ideas from the case that  $\Lambda_1 > 0$ . We use the same  $\tau$  as in the case that  $\Lambda < 0$  but instead of  $\mu$  we take

$$\mu_k = \frac{\log(\sqrt{5}c_\alpha/(\alpha^k - 1))}{\log \alpha}$$

for each possible value of  $n - n_1 = k = 1, 2, \dots, 271$ . Using Lemma 1 with this setting we also obtain in this case that  $n - n_1 \leq 271$  implies  $m - m_1 \leq 219$ .

Now let us turn to the case that  $m - m_1 \leq 212$  and let us consider inequality (9). Recall that

$$\Lambda_2 = n \log \alpha - m_1 \log \alpha_T + \log \left( \frac{1}{\sqrt{5}c_\alpha(\alpha_T^{m-m_1} - 1)} \right)$$

and let us assume that  $n - n_1 \geq 20$ . Then we have

$$|\Lambda_2| < \frac{4.44\alpha^4}{\alpha^{n-n_1}}.$$

Assuming that  $\Lambda_2 > 0$ , we get

$$0 < n \left( \frac{\log \alpha}{\log \alpha_T} \right) - m_1 + \frac{\log(1/(\sqrt{5}c_\alpha(\alpha_T^{m-m_1} - 1)))}{\log \alpha_T} < \frac{4.44\alpha^4}{(\log \alpha_T)\alpha^{n-n_1}} < 50\alpha^{-(n-n_1)}.$$

Once again we apply Lemma 1 with the same  $\tau$  and  $M$  as for the case  $\Lambda > 0$  before. We take  $(A, B) = (50, \alpha)$  and

$$\mu_k = \frac{\log(1/(\sqrt{5}c_\alpha(\alpha_T^k - 1)))}{\log \alpha_T}$$

for every possible value  $m - m_1 = k = 1, 2, \dots, 212$ . If we use again the 104-th convergent of  $\tau$ , i.e. we put  $q = q_{104}$ , then for each  $k$  that yields a positive  $\varepsilon$ , we get  $\varepsilon > 0.0012$ . Therefore we get

$$n - n_1 < \frac{\log(50q_{104}/0.0012)}{\log \alpha} < 280$$

in these cases. But for  $k = 90$  we get a negative  $\varepsilon$ . In this case we consider the 105-th convergent  $\frac{p}{q} = \frac{p_{105}}{q_{105}}$  of  $\tau$  instead. Let us note that

$$q_{105} = 20120013979896675119357414743592977629715414121119669783.$$

Now we obtain in the case  $k = 90$  that  $\varepsilon > 0.46$ . Thus

$$n - n_1 < \frac{\log(50q_{105}/0.46)}{\log \alpha} < 275.$$

In the case that  $\Lambda_2 < 0$  we follow again the ideas from the case that  $\Lambda_2 > 0$ . Of course we choose

$$\tau = \frac{\log \alpha_T}{\log \alpha} \quad \text{and} \quad \mu_k = \frac{\log(\sqrt{5}c_\alpha(\alpha_T^k - 1))}{\log \alpha}.$$

Applying Lemma 1 for all possible values of  $m - m_1 = k = 1, \dots, 212$  also yields in this case that  $n - n_1 \leq 279$ .

Let us summarize the above computations. First we got that either  $n - n_1 \leq 271$ , or  $m - m_1 \leq 212$ . If we assume that  $n - n_1 \leq 271$ , then we deduce that  $m - m_1 \leq 219$ , and if we assume that  $m - m_1 \leq 212$ , then we deduce that  $n - n_1 \leq 279$ . Altogether we obtain  $n - n_1 \leq 279$  and  $m - m_1 \leq 219$ .

For the last step in our reduction process we consider inequality (10). Recall that

$$\Lambda_3 = n_1 \log \alpha - m_1 \log \alpha_T + \log \left( \frac{\alpha^{n-n_1} - 1}{\sqrt{5}c_\alpha(\alpha_T^{m-m_1} - 1)} \right).$$

Since we assume that  $n \geq 300$ , inequality (10) implies that

$$|\Lambda_3| < \frac{3.28\alpha^4}{\alpha^n}.$$

Let us assume that  $\Lambda_3 > 0$ . Then

$$0 < n_1 \left( \frac{\log \alpha}{\log \alpha_T} \right) - m_1 + \frac{\log((\alpha^k - 1)/(\sqrt{5}c_\alpha(\alpha_T^l - 1)))}{\log \alpha_T} < \frac{3.28\alpha^4}{(\log \alpha_T)\alpha^n} < 37\alpha^{-n},$$

where  $(k, l) = (n - n_1, m - m_1)$ . We apply Lemma 1 once more with the same  $\tau$  and  $M$  as for the case when  $\Lambda > 0$ . Moreover, we take  $(A, B) = (37, \alpha)$ , and put

$$\mu_{k,l} = \frac{\log((\alpha^k - 1)/(\sqrt{5}c_\alpha(\alpha_T^l - 1)))}{\log \alpha_T}$$

for  $1 \leq k \leq 279$  and  $1 \leq l \leq 219$ . We consider the 104-th convergent  $\frac{p}{q} = \frac{p_{104}}{q_{104}}$ . For all pairs  $(k, l)$  such that  $\varepsilon$  is positive we have indeed  $\varepsilon > 2.8 \times 10^{-6}$ . Thus for these pairs  $(k, l)$  Lemma 1 yields that

$$n \leq \frac{\log(37q_{104}/0.0000028)}{\log \alpha} < 292.$$

For all the remaining pairs  $(k, l)$  which yield a negative  $\varepsilon$ , we consider the 105-th convergent  $\frac{p}{q} = \frac{p_{105}}{q_{105}}$  instead. And for all those pairs  $(k, l)$  the quantity  $\varepsilon$  is positive for this choice of  $q$ . In particular, we have that  $\varepsilon > 0.0018$  for all these cases, hence

$$n \leq \frac{\log(37q_{105}/0.0018)}{\log \alpha} < 286.$$

In the case that  $A_3 < 0$  the method is similar. In particular we have to apply Lemma 1 with

$$\tau = \frac{\log \alpha_T}{\log \alpha} \quad \text{and} \quad \mu_{k,l} = \frac{\log((\sqrt{5} c_\alpha (\alpha_T^l - 1))/(\alpha^k - 1))}{\log \alpha}.$$

However, we obtain in this case the slightly smaller bound  $n < 289$ .

Altogether our reduction procedure yields the upper bound  $n \leq 291$ . However, this contradicts our assumption that  $n \geq 300$ . Thus Theorem 1 is proved.

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